OVP Guide to Using Processor Models

Model Specific Information for variant ARM_ARM926EJ-S

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Model Release Status

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Table of Contents

1 Overview

This document provides the details of an OVP Fast Processor Model variant.

OVP Fast Processor Models are written in C and provide a C API for use in C based platforms. The models also provide a native interface for use in SystemC TLM2 platforms.

The models are written using the OVP VMI API that provides a Virtual Machine Interface that defines the behavior of the processor. The VMI API makes a clear line between model and simulator allowing very good optimization and world class high speed performance.

Most models are provided as a binary shared object and also as source. This allows the download and use of the model binary or the use of the source to explore and modify the model.

The models are run through an extensive QA and regression testing process and most model families are validated using technology provided by the processor IP owners.

There is a companion document (OVP Guide to Using Processor Models) which explains the general concepts of OVP Fast Processor Models and their use. It is downloadable from the OVPworld website documentation pages.

1.1 Description

ARM Processor Model

1.2 Licensing

Usage of binary model under license governing simulator usage.

Note that for models of ARM CPUs the license includes the following terms:

Licensee is granted a non-exclusive, worldwide, non-transferable, revocable licence to:

If no source is being provided to the Licensee: use and copy only (no modifications rights are granted) the model for the sole purpose of designing, developing, analyzing, debugging, testing, verifying, validating and optimizing software which: (a) (i) is for ARM based systems; and (ii) does not incorporate the ARM Models or any part thereof; and (b) such ARM Models may not be used to emulate an ARM based system to run application software in a production or live environment.

If source code is being provided to the Licensee: use, copy and modify the model for the sole purpose of designing, developing, analyzing, debugging, testing, verifying, validating and optimizing software which: (a) (i) is for ARM based systems; and (ii) does not incorporate the ARM Models or any part thereof; and (b) such ARM Models may not be used to emulate an ARM based system to run application software in a production or live environment.

In the case of any Licensee who is either or both an academic or educational institution the purposes shall be limited to internal use.

Except to the extent that such activity is permitted by applicable law, Licensee shall not reverse engineer, decompile, or disassemble this model. If this model was provided to Licensee in Europe, Licensee shall not reverse engineer, decompile or disassemble the Model for the purposes of error correction.

The License agreement does not entitle Licensee to manufacture in silicon any product based on this model.

The License agreement does not entitle Licensee to use this model for evaluating the validity of any ARM patent.

Source of model available under separate Imperas Software License Agreement.

1.3 Limitations

Instruction pipelines are not modeled in any way. All instructions are assumed to complete immediately. This means that instruction barrier instructions (e.g. ISB, CP15ISB) are treated as NOPs, with the exception of any undefined instruction behavior, which is modeled. The model does not implement speculative fetch behavior. The branch cache is not modeled.

Caches and write buffers are not modeled in any way. All loads, fetches and stores complete immediately and in order, and are fully synchronous (as if the memory was of Strongly Ordered or Device-nGnRnE type). Data barrier instructions (e.g. DSB, CP15DSB) are treated as NOPs, with the exception of any undefined instruction behavior, which is modeled. Cache manipulation instructions are implemented as NOPs, with the exception of any undefined instruction behavior, which is modeled.

Real-world timing effects are not modeled: all instructions are assumed to complete in a single cycle.

TLBs are architecturally-accurate but not device accurate. This means that all TLB maintenance and address translation operations are fully implemented but the cache is larger than in the real device.

1.4 Verification

Models have been extensively tested by Imperas. ARM9 models have been successfully used by customers to simulate Linux and Nucleus on ArmIntegrator virtual platforms.

1.5 Features

1.5.1 Core Features

Thumb instructions are supported.

Trivial Jazelle extension is implemented.

1.5.2 Memory System

FCSE extension is implemented.

VMSA address translation is implemented.

1 ITCM is implemented.

1 DTCM is implemented.

1.6 Integration Support

This model implements a number of non-architectural pseudo-registers and other features to facilitate integration.

1.6.1 Memory Transaction Query

Two registers are intended for use within memory callback functions to provide additional information about the current memory access. Register transactPL indicates the processor execution level of the current access (0-3). Note that for load/store translate instructions (e.g. LDRT, STRT) the reported execution level will be 0, indicating an EL0 access. Register transactAT indicates the type of memory access: 0 for a normal read or write; and 1 for a physical access resulting from a page table walk.

1.6.2 Page Table Walk Query

A banked set of registers provides information about the most recently completed page table walk. There are up to six banks of registers: bank 0 is for stage 1 walks, bank 1 is for stage

2 walks, and banks 2-5 are for stage 2 walks initiated by stage 1 level 0-3 entry lookups, respectively. Banks 1-5 are present only for processors with virtualization extensions. The currently active bank can be set using register PTWBankSelect. Register PTWBankValid is a bitmask indicating which banks contain valid data: for example, the value 0xb indicates that banks 0, 1 and 3 contain valid data.

Within each bank, there are registers that record addresses and values read during that page table walk. Register PTWBase records the table base address. Registers PTWAddressL0-PTWAddressL3 record the addresses of level 0 to level 3 entries read, respectively, and register PTWAddressValid is a bitmask indicating which address registers contain valid data: for example, the value 0xe indicates that PTWAddressL1-PTWAddressL3 are valid but PTWAddressL0 is not. Registers PTWValueL0-PTWValueL3 contain entry values read at level 0 to level 3. Register PTWInput contains the input address that starts a walk and Register PTWOutput contains the result address (valid only if the page table walk completes). Register PTWValueValid is a bitmask indicating which value registers contain valid data: bits 0-3 indicate PTWValueL0-PTWValueL3, respectively, bit 4 indicates PTWBase, bit 5 indicates PTWInput and bit 6 indicates PTWOutput.

1.6.3 Artifact Page Table Walks

Registers are also available to enable a simulation environment to initiate an artifact page table walk (for example, to determine the ultimate PA corresponding to a given VA). Register PTWI_EL1S initiates a secure EL1 table walk for a fetch. Register PTWD_EL1S initiates a secure EL1 table walk for a load or store (note that current ARM processors have unified TLBs, so these registers are synonymous). Registers PTW[ID]_EL1NS initiate walks for non-secure EL1 accesses. Registers PTW[ID]_EL2 initiate EL2 walks. Registers PTW[ID]_S2 initiate stage 2 walks. Registers PTW[ID]_EL3 initiate AArch64 EL3 walks. Finally, registers PTW[ID]_current initiate current-mode walks (useful in a memory callback context). Each walk fills the query registers described above.

1.6.4 MMU and Page Table Walk Events

Two events are available that allow a simulation environment to be notified on MMU and page table walk actions. Event mmuEnable triggers when any MMU is enabled or disabled. Event pageTableWalk triggers on completion of any page table walk (including artifact walks).

1.6.5 Artifact Address Translations

A simulation environment can trigger an artifact address translation operation by writing to the architectural address translation registers (e.g. ATS1CPR). The results of such translations are written to an integration support register artifactPAR, instead of the architectural PAR register. This means that such artifact writes will not perturb architectural state.

1.6.6 Halt Reason Introspection

An artifact register HaltReason can be read to determine the reason or reasons that a processor is halted. This register is a bitfield, with the following encoding: bit 0 indicates the processor has executed a wait-for-event (WFE) instruction; bit 1 indicates the processor has executed a wait-for-interrupt (WFI) instruction; and bit 2 indicates the processor is held in reset.

1.6.7 System Register Access Monitor

If parameter enableSystemMonitorBus is True, an artifact 32-bit bus "SystemMonitor" is enabled for each PE. Every system register read or write by that PE is then visible

as a read or write on this artifact bus, and can therefore be monitored using callbacks installed in the client environment (use opBusReadMonitorAdd/opBusWriteMonitorAdd or icmAddBusReadCallback/icmAddBusWriteCallback, depending on the client API). The format of the address on the bus is as follows:

bits 31:26 - zero

bit 25 - 1 if AArch64 access, 0 if AArch32 access

bit 24 - 1 if non-secure access, 0 if secure access

bits 23:20 - CRm value

bits 19:16 - CRn value

bits 15:12 - op2 value

bits 11:8 - op1 value

bits 7:4 - op0 value (AArch64) or coprocessor number (AArch32)

bits 3:0 - zero

As an example, to view non-secure writes to writes to CNTFRQ_EL0 in AArch64 state, install a write monitor on address range 0x020e0330:0x020e0333.

1.6.8 System Register Implementation

If parameter enableSystemBus is True, an artifact 32-bit bus "System" is enabled for each PE. Slave callbacks installed on this bus can be used to implement modified system register behavior (use opBusSlaveNew or icmMapExternalMemory, depending on the client API). The format of the address on the bus is the same as for the system monitor bus, described above.

2 Configuration

2.1 Location

The model source and object file is found in the VLNV tree at:

arm.ovpworld.org/processor/arm/1.0

2.2 GDB Path

The default GDB for this model is found at:

$IMPERAS_HOME/lib/$IMPERAS_ARCH/gdb/arm-none-eabi-gdb

2.3 Semi-Host Library

The default semi-host library file is found in the VLNV tree at :

arm.ovpworld.org/semihosting/armNewlib/1.0

2.4 Processor Endian-ness

This model can be set to either endian-ness (normally by a pin, or the ELF code).

2.5 QuantumLeap Support

This processor is qualified to run in a QuantumLeap enabled simulator.

2.6 Processor ELF Code

The ELF code supported by this model is: 0x28

3 Other Variants in this Model

Table 1.

Variant
ARMv4T
ARMv4xM
ARMv4
ARMv4TxM
ARMv5xM
ARMv5
ARMv5TxM
ARMv5T
ARMv5TExP
ARMv5TE
ARMv5TEJ
ARMv6
ARMv6K
ARMv6T2
ARMv6KZ
ARMv7
ARM7TDMI
ARM7EJ-S
ARM720T
ARM920T
ARM922T
ARM926EJ-S
ARM940T
ARM946E
ARM966E
ARM968E-S
ARM1020E
ARM1022E
ARM1026EJ-S
ARM1136J-S
ARM1156T2-S
ARM1176JZ-S
Cortex-R4
Cortex-R4F
Cortex-A5UP
Cortex-A5MPx1
Cortex-A5MPx2
Cortex-A5MPx3

Cortex-A5MPx4
Cortex-A8
Cortex-A9UP
Cortex-A9MPx1
Cortex-A9MPx2
Cortex-A9MPx3
Cortex-A9MPx4
Cortex-A7UP
Cortex-A7MPx1
Cortex-A7MPx2
Cortex-A7MPx3
Cortex-A7MPx4
Cortex-A15UP
Cortex-A15MPx1
Cortex-A15MPx2
Cortex-A15MPx3
Cortex-A15MPx4
Cortex-A17MPx1
Cortex-A17MPx2
Cortex-A17MPx3
Cortex-A17MPx4
AArch32
AArch64
Cortex-A53MPx1
Cortex-A53MPx2
Cortex-A53MPx3
Cortex-A53MPx4
Cortex-A57MPx1
Cortex-A57MPx2
Cortex-A57MPx3
Cortex-A57MPx4
Cortex-A72MPx1
Cortex-A72MPx2
Cortex-A72MPx3
Cortex-A72MPx4
MultiCluster

4 Bus Ports

Table 2.

Type	Name	min	max	Description
master (initiator)	ITCM0	32	32	instruction TCM

master (initiator)	DTCM0	32	32	data TCM
master (initiator)	INSTRUCTION	32	32	
master (initiator)	DATA	32	32	

5 Net Ports

Table 3.

Name	Type	Description
reset	input	Processor reset, active high
fiq	input	FIQ interrupt, active high (negation of nFIQ)
irq	input	IRQ interrupt, active high (negation of nIRQ)
sei	input	System error interrupt, active high

6 FIFO Ports

No FIFO Ports in this model.

7 Parameters

Table 4.

Name	Type	Description
verbose	Boolean	Specify verbosity of output
showHiddenRegs	Boolean	Show hidden registers during register tracing
UAL	Boolean	Disassemble using UAL syntax
enableSystemBus	Boolean	Add 32-bit artifact System bus port, allowing system registers to be externally implemented
enableSystemMonitorBus	Boolean	Add 32-bit artifact SystemMonitor bus port, allowing system register accesses to be externally monitored
compatibility	Enumeration	Specify compatibility mode ISA=0 gdb=1 nopSVC=2
override_debugMask	Uns32	Specifies debug mask, enabling debug output for model components
override_fcsePresent	Boolean	Specifies that FCSE is present (if true)
override_SCTLR_V	Boolean	Override SCTLR.V with the passed value (enables high vectors)
override_SCTLR_CP15BEN_Present	Boolean	Enable ARMv7 SCTLR.CP15BEN bit (CP15 barrier enable)
override_MIDR	Uns32	Override MIDR/MIDR_EL1 register
override_CTR	Uns32	Override CTR/CTR_EL0 register
override_TLBTR	Uns32	Override TLBTR register
override_CLIDR	Uns32	Override CLIDR/CLIDR_EL1 register

override_AIDR	Uns32	Override AIDR/AIDR_EL1 register
override_ERG	Uns32	Specifies exclusive reservation granule
override_STRoffsetPC12	Boolean	Specifies that STR/STR of PC should do so with 12:byte offset from the current instruction (if true), otherwise an 8:byte offset is used
override_fcseRequiresMMU	Boolean	Specifies that FCSE is active only when MMU is enabled (if true)
override_ignoreBadCp15	Boolean	Specifies whether invalid coprocessor 15 access should be ignored (if true) or cause Invalid Instruction exceptions (if false)
override_SGIDisable	Boolean	Override whether GIC SGIs may be disabled (if true) or are permanently enabled (if false)
override_condUndefined	Boolean	Force undefined instructions to take Undefined Instruction exception even if they are conditional
override_deviceStrongAligned	Boolean	Force accesses to Device and Strongly Ordered regions to be aligned
override_Control_V	Boolean	Override SCTLR.V with the passed value (deprecated, use override_SCTLR_V)
override_MainId	Uns32	Override MIDR register (deprecated, use override_MIDR)
override_CacheType	Uns32	Override CTR register (deprecated, use override_CTR)
override_TLBType	Uns32	Override TLBTR register (deprecated, use override_TLBTR)

8 Execution Modes

Table 5.

Name	Code
User	16
FIQ	17
IRQ	18
Supervisor	19
Abort	23
Undefined	27
System	31

9 Exceptions

Table 6.

Name	Code
Reset	0

Undefined	1
SupervisorCall	2
PrefetchAbort	5
DataAbort	6
IRQ	8
FIQ	9

10 Hierarchy of the model

A CPU core may allow the user to configure it to instance many processors of a Symmetrical Multi Processor (SMP). A CPU core may also have sub elements within a processor, for example hardware threading blocks.
OVP processor models can be written to include SMP blocks and to have many levels of hierarchy.

Some OVP CPU models may have a fixed hierarchy, and some may be configured by settings in a configuration register. Please see the register definitions of this model.

This model documentation shows the settings and hierarchy of the default settings for this model variant.

10.1 Level 1: CPU

This level in the model hierarchy has 5 commands.

This level in the model hierarchy has 10 register groups:

Table 7.

Group name	Registers
Core	16
Control	3
User	7
FIQ	8
IRQ	3
Supervisor	3
Undefined	3
Abort	3
Coprocessor_32_bit	42
Integration_support	22

This level in the model hierarchy has no children.

11 Model Commands

11.1 Level 1: CPU

Table 8.

Name	Arguments
debugflags	
dumpTLB	[-all]
isync	specify instruction address range for synchronous execution
itrace	enable or disable instruction tracing
validateTLB	[-all] [-verbose]

12 Registers

12.1 Level 1: CPU

12.1.1 Core

Table 9.

Name	Bits	Initial value (Hex)		Description
r0	32	0	rw	
r1	32	0	rw	
r2	32	0	rw	
r3	32	0	rw	
r4	32	0	rw	
r5	32	0	rw	
r6	32	0	rw	
r7	32	0	rw	
r8	32	0	rw	
r9	32	0	rw	
r10	32	0	rw	
r11	32	0	rw	frame pointer
r12	32	0	rw	
sp	32	0	rw	stack pointer
lr	32	0	rw	
pc	32	0	rw	program counter

12.1.2 Control

Table 10.

Name	Bits	Initial value (Hex)		Description

fps	32	0	rw	archaic FPSCR view (for gdb)
cpsr	32	d3	rw	
spsr	32	0	rw	

12.1.3 User

Table 11.

Name	Bits	Initial value (Hex)		Description
r8_usr	32	0	rw	
r9_usr	32	0	rw	
r10_usr	32	0	rw	
r11_usr	32	0	rw	
r12_usr	32	0	rw	
sp_usr	32	0	rw	
lr_usr	32	0	rw	

12.1.4 FIQ

Table 12.

Name	Bits	Initial value (Hex)		Description
r8_fiq	32	0	rw	
r9_fiq	32	0	rw	
r10_fiq	32	0	rw	
r11_fiq	32	0	rw	
r12_fiq	32	0	rw	
sp_fiq	32	0	rw	
lr_fiq	32	0	rw	
spsr_fiq	32	0	rw	

12.1.5 IRQ

Table 13.

Name	Bits	Initial value (Hex)		Description
sp_irq	32	0	rw	
lr_irq	32	0	rw	
spsr_irq	32	0	rw	

12.1.6 Supervisor

Table 14.

Name	Bits	Initial value (Hex)		Description
sp_svc	32	0	rw	
lr_svc	32	0	rw	
spsr_svc	32	0	rw	

12.1.7 Undefined

Table 15.

Name	Bits	Initial value (Hex)		Description
sp_undef	32	0	rw	
lr_undef	32	0	rw	
spsr_undef	32	0	rw	

12.1.8 Abort

Table 16.

Name	Bits	Initial value (Hex)		Description
sp_abt	32	0	rw	
lr_abt	32	0	rw	
spsr_abt	32	0	rw	

12.1.9 Coprocessor_32_bit

Table 17.

Name	Bits	Initial value (Hex)		Description
CONTEXTIDR	32	0	rw	Context ID
CTR	32	1d112152	r-	Cache Type
CleanDCacheLineMVA	32	-	-w	Data Cache Line Clean by VA
CleanDCacheLineSW	32	-	-w	Data Cache Line Clean by Set/Way
CleanInvalDCacheLineMVA	32	-	-w	Data Cache Line Clean and Invalidate by VA
CleanInvalDCacheLineSW	32	-	-w	Data Cache Line Clean and Invalidate by Set/Way
DACR	32	0	rw	Domain Access Control
DCLR	32	fff0	rw	Data Cache Lockdown
DFSR	32	0	rw	Data Fault Status
DTCMRR	32	18	rw	DTCM Region
DTLBIALL	32	-	-w	Invalidate Entire Data TLB

DTLBIMVA	32	-	-w	Invalidate Data TLB by VA
DTLBLR	32	0	rw	TLB Lockdown
DataSynchBarrier	32	-	-w	Data Synchronization Barrier
FAR	32	0	rw	Fault Address
FCSEIDR	32	0	rw	FCSE Process ID
ICLR	32	fff0	rw	Instruction Cache Lockdown
IFSR	32	0	rw	Instruction Fault Status
ITCMRR	32	18	rw	ITCM Region
ITLBIALL	32	-	-w	Invalidate Entire Instruction TLB
ITLBIMVA	32	-	-w	Invalidate Instruction TLB by VA
InvalDCache	32	-	-w	Invalidate Data Cache
InvalDCacheLineMVA	32	-	-w	Invalidate Data Cache Line by VA
InvalDCacheLineSW	32	-	-w	Invalidate Data Cache Line by Set/Way
InvalICache	32	-	-w	Invalidate Instruction Cache
InvalICacheLineMVA	32	-	-w	Invalidate Instruction Cache Line by VA
InvalICacheLineSW	32	-	-w	Invalidate Instruction Cache Line by Set/Way
InvalUnified	32	-	-w	Invalidate Unified Cache
JIDR	32	0	rw	Jazelle ID
JMCR	32	0	rw	Jazelle Main Configuration
JOSCR	32	0	rw	Jazelle OS Control
MIDR	32	41069265	r-	Main ID
PrefetchICacheLine	32	-	-w	Prefetch Instruction Cache Line
SCTLR	32	50078	rw	System Control
TCMTR	32	10001	r-	TCM Type
TLBIALL	32	-	-w	Invalidate Entire Unified TLB
TLBIMVA	32	-	-w	Invalidate Unified TLB by VA
TTBR	32	0	rw	Translation Table Base
TestCleanDCache	32	400000d3	r-	Data Cache Test and Clean
TestCleanInvalDCache	32	400000d3	r-	Data Cache Test, Clean and Invalidate
WaitForInterrupt	32	-	-w	Wait For Interrupt
WaitForInterrupt2	32	-	-w	Wait For Interrupt

12.1.10 Integration_support

Table 18.

Name	Bits	Initial value (Hex)		Description
transactPL	32	1	r-	privilege level of current memory transaction
transactAT	32	0	r-	current memory transaction type: PA=1, VA=0
artifactPAR	64	0	r-	result of address translation for artifact write to ATS1CPR etc
PTWBankValid	8	0	r-	bitmask of valid banks (0x01 is stage 1, 0x02 is stage 2, 0x04-0x20 are stage 2 walks

				initiated by stage 1 level 0-3 entry lookups, respectively)
PTWAddressValid	8	0	r-	bitmask of valid bits for each of PTWAddressL0...PTWAddressL3, PTWBase, PTWInput and PTWOutput in current bank
PTWValueValid	8	0	r-	bitmask of valid bits for each of PTWValueL0...PTWValueL3 in current bank
PTWAddressL0	64	0	r-	current bank PTW address, level 0
PTWAddressL1	64	0	r-	current bank PTW address, level 1
PTWAddressL2	64	0	r-	current bank PTW address, level 2
PTWAddressL3	64	0	r-	current bank PTW address, level 3
PTWValueL0	64	0	r-	current bank PTW value, level 0
PTWValueL1	64	0	r-	current bank PTW value, level 1
PTWValueL2	64	0	r-	current bank PTW value, level 2
PTWValueL3	64	0	r-	current bank PTW value, level 3
PTWBase	64	0	r-	current bank PTW table base address
PTWInput	64	0	r-	current bank PTW input address
PTWOutput	64	0	r-	current bank PTW output address
PTWI_EL1NS	64	-	-w	perform EL1(NS) stage 1 page table walk for fetch, filling PTW query registers
PTWD_EL1NS	64	-	-w	perform EL1(NS) stage 1 page table walk for load/store, filling PTW query registers
PTWI_current	64	-	-w	perform current mode page table walk for fetch, filling PTW query registers
PTWD_current	64	-	-w	perform current mode page table walk for load/store, filling PTW query registers
HaltReason	8	0	r-	bit field indicating halt reason

#

